Technologies and Models for En Route Commerce

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Definitions

• Electronic Commerce:

Buying and selling goods using electronic transaction processing technologies (that require little or no intervention on the part of the buyer or seller).

• En Route Commerce:

Electronic commerce conducted while traveling from one location to another.



Enroute Commerce for Personal Travel

• General Travel

Lodging Pricing/discounts and availability Reservations Electronic check-in/keys Restaurants and Entertainment Information Reservations Electronic ticketing

• Automobile Travel

Gasoline pricing/purchases Parking pricing/reservations/auctions Roadside service



Enroute Commerce for Personal Travel (cont.)

• Air/Bus/Train Travel

Airport Services Travel "aids" (e.g., movies) Porters Purchase rail tickets/passes/e-tickets

• Taxi Travel

Competitive, dynamic pricing

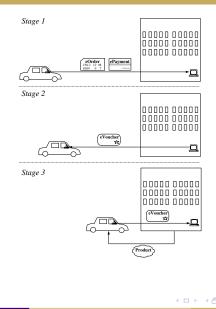


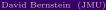
"Commercial" Enroute Commerce

- Mobile asset management by carrier
 - Load matching Driver utilization Equipment utilization Dynamic re-routing
- Mobile asset management by shipper Dynamic redployment of shipments Improved responsiveness (info and controls)
- Personal services (as above)



Type-1 Transactions

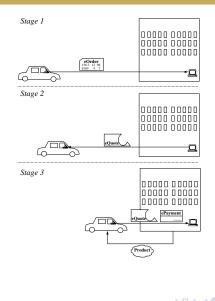




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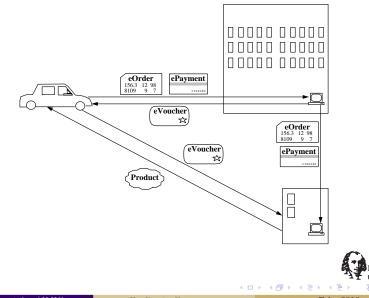
Type-2 Transactions



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AMES IADISON NIVERSITY® En Route Commerce Contractual Exchanges

Type-3 Transactions (Using A Central Office)

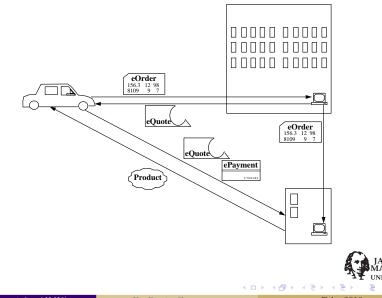


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En Route Commerce Contractual Exchanges

Type-4 Transactions (Using A Central Office)



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Taxi Services: Overview

• The Past:

Manual requests/reservations Drop fee and per mile (e.g., NYC) or zone-based fares (e.g., DC) Cash payments

- Today/The Near Future: Electronic requests/reservations Electronic payment
- The Future:

Real-time pricing?



Taxi Pricing: Assumptions

- Short, Frequent Auctions: Every 5 minutes
- Falling Price Auction:

Firms announce prices for particular trips Individuals accept/reject Firms adjust prices

All individuals receive "final" price



Modeling Assumptions

- Continuum of utility-maximizing passengers
- Continuum of homogeneous vehicles
- An available alternative (i.e., elastic demand)
- One time period/auction
- Small number of (or one) profit-maximizing firms
- One or more OD-pairs



- A B D

Model

• Related Models:

Spatial Price Equilibrium (Samuleson, 1952; Takayama and Judge, 1971)

Spatial Oligopoly (Weskamp, 1985; Harker, 1986; Nagurney, 1987)

Spatial Oligopoly with Arbitrageurs (Friesz and Bernstein, 1992)

• Important Difference:

Customers can move

"Product" can only be moved by the firm



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Formulations

• Complementarity Problem:

Given a function $G: \mathbb{R}^n_+ \to \mathbb{R}^n$, find an $x \in \mathbb{R}^n$ such that:

$$\begin{aligned} x_i &\geq 0 \quad G_i(x) \geq 0 \quad G_i(x) \\ x_i \text{ free } \quad G_i(x) &= 0 \quad G_i(x) \\ x_i &= 0 \quad \forall i \in T_2, \end{aligned}$$
(1)

where $T_1 \cup T_2$ is a partition of the indices $\{1, 2, \dots, n\}$.

• Variational Inequality Problem:

Given a function $G: R_+^n \to R^n$ and a (closed convex) set $X \subseteq R^n$, find an x such that:

$$G(x)^T(y-x) \ge 0 \ \forall y \in X.$$
(2)

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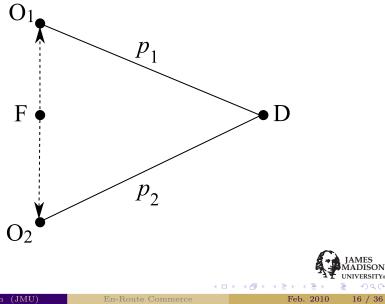
Existence and Uniqueness

- Existence Conditions: Involve continuity
- Uniqueness Conditions:

Involve some form of monotonicity



A Monopoly Example



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A Monopoly Example (cont.)

• Demand:

$$p_1 = 4 - q_1$$

$$p_2 = 6 - 2q_2$$

• Marginal Cost:

C' = 0.2 for all trips



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A Monopoly Example (cont.)

• Prohibit
vely High Walking Cost: $p_1 = 2.10, \, p_2 = 3.10$

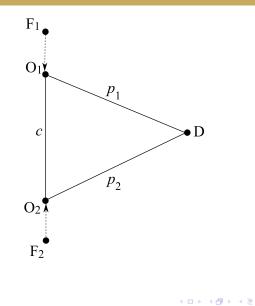
$$q_1 = 1.90, q_2 = 1.45$$

$$\pi_1 = 7.82$$

• Walking Cost of 0.5: $p_1 = 2.27, p_2 = 2.27$ $q_1 = 1.73, q_2 = 1.62$ $\pi_2 = 7.73$



An Oligopoly Example





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An Oligopoly Example (cont.)

• Demand:

$$p_1 = 4 - q_1$$

$$p_2 = 13 - 2q_2$$

- Costs:
 - $C'_{11} = 0.2$ $C'_{12} = 0.2 + c$ $C'_{21} = 2.1 + c$ $C'_{22} = 2.1$ where c = 3.0



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An Oligopoly Example (cont.)

• Prohibitively High Walking Cost: $p_1 = 2.10, p_2 = 6.10$

$$q_{11} = 1.90, q_{12} = 1.45, \pi = 7.82$$

$$q_{21} = 0.00, \, q_{22} = 2.00, \, \pi = 8.00$$

• Low Walking Cost:

Firms can "allow" walkers

Firms can "prevent" walkers (by pricing appropriately)



An Oligopoly Example (cont.)

• Assume:

Walking cost of 3.0 (i.e., equal to c)

Firms prevent walking

• One Solution:

 $p_1 = 2.40 \uparrow, p_2 = 5.40 \downarrow$

$$q_{11} = 1.60 \downarrow, q_{12} = 1.80 \uparrow$$

 $q_{21} = 0.00, \, q_{22} = 2.00$



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En Route Gasoline Purchases: Overview

• The Past:

Historical information about prices Cash payments

• Today/The Near Future:

Location of some facilities Real-time reported prices at some facilities Electronic payment

• The Future:

Comprehensive real-time guaranteed prices Decision support system that incorporates gasoline purchases



Assumptions

- 1. Given an origin and destination, the decision support system determines the best gas station and route.
- 2. There is only one gas station at any node. (We can still handle real-world situations by adding "dummy nodes".)
- 3. We only consider *short trips* (i.e., the amount of fuel in the vehicle is sufficient to reach all gas stations and the tank holds enough fuel to get to the destination).
- 4. The driver only wants to stop (for gasoline) once.
- 5. The driver's value of time is a known constant (α).
- 6. The driver always drives at exactly the speed limit.



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Notation

 ${\cal O}$ denotes the origin and ${\cal D}$ denotes the destination

 $a_{ij} = 1$ if link j is directed out of node i, $a_{ij} = -1$ if link j is directed into node i, and $a_{ij} = 0$ otherwise

$$b_O = 1, b_D = -1, \text{ and } b_i = 0, i \in \mathcal{N} - \{O, D\}$$

 p_g denotes the price (per gallon) of gasoline at station g

m denotes the *fuel efficiency* (in miles per gallon) of the vehicle

 s_{ij} denotes the speed limit on link (i, j)

 c_{ij} denotes the total cost of traveling on link (i, j)



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An Observation

• The Observation:

We can ignore the gasoline in the vehicle at the start of the problem

• Different Rationales:

It's a fixed cost

Opportunity cost of the gas in the tank is the price you purchase gas at today



Link Costs

$$c_{ij} = p \cdot (d_{ij}/m) + \alpha \cdot (d_{ij}/s_{ij}) \tag{3}$$

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An Algorithmic Approach

Use a multi-copy network in which there is a copy of the physical link for each gasoline price

Find the shortest path from O to D via the gas statione (i.e., through an intermediate node)



The Shortest Path (S.P.) Problem

$$\begin{array}{rcl} \min_{x} & c^{\top}x \\ \text{s.t.} & Ax &= b \\ & x &\in \{0,1\}^n \end{array}$$



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(4)

S.P. Through an Intermediate Node

Let h_I denote an *n*-vector in which all elements that that correspond to arcs with head node I (i.e., that are inbound to node I) have a value of 1 and all other elements have a value of 0.

$$\begin{array}{lll} \min_{x} & c^{\top}x \\ \text{s.t.} & Ax &= b \\ & x &\in \{0,1\}^{n} \\ & h_{I}^{\top}x &>= 1 \end{array}$$

$$(5)$$



S.P. Through an Intermediate Node (cont.)

• A Well-Known Algorithm:

Find the shortest path from ${\cal O}$ to ${\cal I}$ and the shortest path from ${\cal I}$ to D

• A Less-Well-Known Algorithm:

Use a label setting algorithm and solve from I to O and D simultaneously (perhaps using two labels if $d_{ij} \neq d_{ji}$ for some i, j)



S.P. Through an Intermediate Node (cont.)

Lagrangean Relaxation

$$\min_{x} \quad c^{\top}x + \lambda(1 - h_{I}^{\top}x)$$

s.t.
$$Ax = b$$

$$x \in \{0, 1\}^{n}$$
 (6)



S.P. Through an Intermediate Node (cont.)

The Shortcoming of Lagrangean Relaxation

$$\min_{x} \quad (c - \lambda h_{I})^{\top} x + \lambda$$
s.t. $Ax = b$

$$x \in \{0, 1\}^{n}$$

$$(7)$$

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 λ has the effect of reducing the distance on the nodes that are in bound to I

 λ must be so large that the important elements of $(d - \lambda h_I)$ are negative



SP Through One of a Set of Nodes

- Solve the k intermediate node problems and choosing the shortest
- Use k+1 labels on each node, one for the "unconstrained" distance and one for the distance through each of the intermediate nodes



Future Research on Real-Time Taxi Pricing

- Walking at the destination
- Taxi transfers (at both ends)
- Supply constraints
- Heterogeneous vehicles
- Multiple periods



Future Research on Side Trips

- Gasoline can be purchased on either the outbound or return trip
- The amount of fuel in the vehicle at the start of the trip constrains the choices
- More than one stop for fuel (and other things) must be made
- Other objectives (e.g., not diverting too far)
- Parallel processing (e.g., shared memory, vector processing)



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